

# Quantum Walks in Discrete and Continuous Time 

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NGT

## Outline

- Quantum Walks
- DTQW \& CTQW
- Quantum Algorithms
- Searching, Mixing, Hitting, \& Graph Traversal
- Connecting the quantum walks
- Relativistic Limits \& Other Approaches
- Applications Present and Future
- Algorithms, Implementations, Quantum State Transfer


## Classical Random Walk

- Consider a process in which a particle can move either left or right, in one dimension, based on an internal state (coin).
- Before each step, the coin is flipped.
- This generates a stochastic process, governed by the


Space following equations for the probability:

$$
\begin{aligned}
& p_{R}(n, t+1)=\frac{1}{2} p_{R}(n-1, t)+\frac{1}{2} p_{L}(n-1, t) \\
& p_{L}(n, t+1)=\frac{1}{2} p_{L}(n+1, t)+\frac{1}{2} p_{R}(n+1, t)
\end{aligned}
$$

## Quantum Random Walk

- Change probabilities to probability amplitudes.
- Change coin flip to unitary transformation.
- E.g. Hadamard Walk:


$$
\begin{aligned}
& \psi_{R}(n, t+1)=\frac{1}{\sqrt{2}} \psi_{R}(n-1, t)+\frac{1}{\sqrt{2}} \psi_{L}(n-1, t) \\
& \psi_{L}(n, t+1)=\frac{1}{\sqrt{2}} \psi_{R}(n+1, t)-\frac{1}{\sqrt{2}} \psi_{L}(n+1, t)
\end{aligned}
$$

Aharonov, Davidovich, and Zagury (1993) Interference!

## Quantum Walk vs. Classical Walk

QW Probabilities
CW Probabilities


Space

$$
\begin{array}{ll}
\psi_{R}(n, t=0)=\frac{1}{\sqrt{2}} \delta_{n, 0} & p_{R}(n, t=0)=\frac{1}{2} \delta_{n, 0} \\
\psi_{L}(n, t=0)=\frac{1}{\sqrt{2}} i \delta_{n, 0} & p_{L}(n, t=0)=\frac{1}{2} \delta_{n, 0}
\end{array}
$$

Ensures symmetric dynamics

## Quantum Interference = Quadratic Speed-up



$$
\Delta x^{2} \sim \Delta t^{2}
$$


$\Delta x^{2} \sim \Delta t$

Interesting Probability Distribution
constant for $-\mathrm{T} / \sqrt{ } 2<\mathrm{x}<\mathrm{T} / \sqrt{ } 2$
Wave equation w/ Dispersion (Knight, Roldan, \& Sipe 2003)

## DTQWs on General Graphs

Discrete-Time Quantum Walk
Use d-dimensional coin operator C for vertices of degree d

Conditional shift operator $\mathbf{S}$, depends on each vertex

$$
S|c, v\rangle=|c, w(c, v)\rangle
$$

$w(c, v)=$ vertex connected to $v$ along edge $c$


## Unitary Mapping:

$$
\begin{aligned}
& U=S(C \otimes I) \\
& |\Psi(T)\rangle=U^{T}|\Psi(0)\rangle
\end{aligned}
$$

Other definitions possible

## DTQWs on General Graphs

Discrete-Time Quantum Walk
Use d-dimensional coin operator C for vertices of degree d

Conditional shift operator $\mathbf{S}$, depends on each vertex

$$
S|c, v\rangle=|c, w(c, v)\rangle
$$

$w(c, v)=$ vertex connected to $v$ along edge $c$

E.g. $d=3$

$$
S|c=1, v=1\rangle=|c=1, v=2\rangle
$$

$$
S|c=2, v=1\rangle=|c=2, v=9\rangle
$$

$$
S|c=3, v=1\rangle=|c=3, v=18\rangle
$$

## Diffusion on General Graphs

Use Laplacian matrix for graph

$$
L=A-D
$$

A = adjacency matrix
D = degree matrix

$D_{i, i}=d_{i}, d_{i}=$ degree of vertex $i$
$A_{i, j}=\left\{\begin{array}{lc}1 & \text { if }(i, j) \text { are connected } \\ 0 & \text { otherwise }\end{array}\right.$

$$
\begin{aligned}
& D_{i, i}=3 \forall i \\
& A_{1,2}=A_{1,9}=A_{1,18}=1
\end{aligned}
$$

## Diffusion on General Graphs

Use Laplacian matrix for graph

$$
L=A-D
$$

A = adjacency matrix

$$
D=\text { degree matrix }
$$


$p=$ probability vector

$$
\sum_{j=1}^{|G|} p_{j}=1 \quad \frac{d p}{d t}=L p \quad p(t)=e^{L t} p(0)
$$

## CTQW on General Graphs

Continuous-Time Quantum Walk
Use Laplacian matrix as Hamiltonian:

$$
H=A-D
$$

Sometimes $H=A$

$$
\Psi=\left\{\psi_{1}, \psi_{2}, \cdots, \psi_{|G|}\right\}
$$


$=$ state vector (probability amplitudes)

$$
\sum_{j=1}^{|G|}\left|\psi_{j}\right|^{2}=1 \quad i \frac{d \Psi}{d t}=H \Psi \quad \Psi(t)=e^{-i H t} \Psi(0)
$$

## DTQW vs CTQW <br> One-dimensional Line



Jacobi polynomial Airy, Bessel fxn approximations


Bessel function

Dynamics can be very similar!

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## Quantum Walk Search

## Grover Search

Similar results by CTQW on complete graph


$$
H=|\psi\rangle\langle\psi|+|w\rangle\langle w|
$$

Inversion about average Oracle

$$
|\psi\rangle=\frac{1}{\sqrt{N}} \sum_{v=1}^{N}|\nu\rangle
$$

= uniform
superposition


$$
U^{T}|\psi\rangle \rightarrow|w\rangle
$$

$$
e^{-i H T}|\psi\rangle \rightarrow|w\rangle
$$

## QW Local Search

| Authors | Graph | Walk | Time |
| :---: | :---: | :---: | :---: |
| Farhi \& Gutmann (1998) | Complete | CTQW | $N^{1 / 2}$ |
| Shenvi, Kempe, \& Whaley (2003) | Hypercube | DTQW by coin | $N^{1 / 2}$ |
| Aaronson \& Ambainis (2003) | D-dim Lattice | DTQW by coin | $\begin{gathered} \mathrm{N}^{1 / 2}(\log \mathrm{~N})^{3 / 2} \mathrm{D}=2 \\ \mathrm{~N}^{1 / 2} \mathrm{D} \geq 3 \end{gathered}$ |
| Childs \& Goldstone (2003) | D-dim Lattice | CTQW | $\begin{gathered} N D=2 \\ N^{5 / 6} D=3 \\ N^{1 / 2} \log N D=4 \\ N^{1 / 2} D>4 \end{gathered}$ |
| Ambainis, Kempe, \& Rivosh (2004) | D-dim Lattice | DTQW by coin | $\begin{gathered} \mathrm{N}^{1 / 2} \log \mathrm{ND}=2 \\ \mathrm{~N}^{1 / 2} \mathrm{D} \geq 3 \end{gathered}$ |
| Childs \& Goldstone (2004) | D-dim Lattice | CTQW + Spin | $\begin{gathered} \mathrm{N}^{1 / 2} \log N \mathrm{D}=2 \\ \mathrm{~N}^{1 / 2} \mathrm{D} \geq 3 \end{gathered}$ |
| Magniez, Nayak, Roland, \& Santha (2007) | D-dim Lattice | DTQW by reflection | $\begin{gathered} N^{1 / 2}(\log N)^{1 / 2} D=2 \\ N^{1 / 2} D \geq 3 \end{gathered}$ |

## Mixing on Quantum Walks

Mixing time $=$ time when probability is essentially uniform over graph

Quantum walks mix faster (often quadratically) than classical walks---and decoherence helps!

Kendon \& Tregenna (2003)

Hypercube


Moore \& Russell (2001)

$$
\left.P_{v}(t)=\left|\langle v| e^{-i H t}\right| a\right\rangle\left.\right|^{2} \rightarrow \frac{1}{2^{d}} \text { at time } \omega t=\pi / 4_{16}
$$

## Hitting on Quantum Walks

Hitting time $=$ time when probability is large $(>\mathrm{I} / \log |\mathrm{G}|)$


Quantum walks hit exponentially faster than random walks on glued trees and hypercube graphs
(One-shot hitting time; others have been studied by Kempe, Kendon, and Brun)

## Hypercube Quantum Walk

- Each vertex of the ddimensional hypercube can be encoded by a set of $\boldsymbol{d}$ bits (E.g. $x=X_{1} X_{2} X_{3} X_{4}$ ).
- Hamiltonian (~Adjacency Matrix) simply flips each bit!

$$
H=\omega \sum_{j=1}^{d} \sigma_{x}^{(j)}
$$

Wavefunction evolves simply:

$$
|\Psi(t)\rangle=\left(e^{-i \omega t \sigma_{x}}\right)^{\otimes d}|\Psi(0)\rangle
$$

At time $T=\frac{\pi}{2 \omega}$ all bits flip!


## Exponential Speedup:

Quantum states propagate from corner to corner in time independent of the hypercube dimension d!

Not a computational speedup

## Glued Trees Quantum Walk



> QW on graph of size $\sim 2^{\mathrm{d}}$ can be represented by walk on line with $2 \mathrm{~d}+1$ sites

QW hits in time ~ $2 \mathrm{~d}+\mathrm{I}$, while RW takes time $\sim 2^{\text {d }}$

Childs et al. (200I)

## Graph Traversal Problem


-Problem: Starting from the entrance (green), find the exit (red). -Resources: Local queries of an oracle.
-Classical Running Time: polynomial in $N$ (number of nodes)

## Graph Traversal Problem


-Problem: Starting from the entrance (green), find the exit (red). -Resources: Local queries of an oracle (unitary operator).
-Quantum Running Time: polynomial in $\log N$.

## Graph Traversal Problem


defect
Exponential Algorithmic Speedup by Quantum Walk! Childs (2003)

Reduction leads to QW on line with a defect

Significant amplitude is transmitted through!

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## Connecting QWs

DTQW (Discrete-Time QW)
$\psi_{R}(n, t+1)=\cos \theta \psi_{R}(n-1, t)-i \sin \theta \psi_{L}(n-1, t)$
$\psi_{L}(n, t+1)=\cos \theta \psi_{L}(n+1, t)-i \sin \theta \psi_{R}(n+1, t)$
$\theta=$ coin rotation angle
CTQW (Continuous-Time QW)
$i \partial_{t} \psi(n, t)=-\gamma[\psi(n-1, t)-2 \psi(n, t)+\psi(n+1, t)]$

- Can one get one from the other?
- What physics drives these walks?


## Feynman Path Integral

- DTQW = Feynman's Checkerboard

- Feynman's Checkboard
= Dirac Equation
$i \hbar \partial_{t} \Psi=\left(-i \hbar c \vec{\alpha} \cdot \nabla+\beta m c^{2}\right) \Psi$
D. Meyer (I996)


## Feynman's Proof



Figure 8 from "Feynman and the visualization of space-time processes", Silvan S. Schweber, Rev. Mod. Phys. 58, 499 (1986).

## Feynman's Proof



Feynman knows his Bessel
Feynman (I946) functions!

## Feynman's Proof



This is the DTQW!
Feynman (I946)

## Relativistic Walks

- Can we use the Dirac equation to understand the DTQW?
- Yes, if we have some method to generate wave-packet solutions
- Look at wave-packet solutions of Dirac equation!
- (Use Momentum Space!)
- Compare with wave-packet solutions of DTQW
- (Use Momentum Space!)
- CTQW comes for free! FWs, Phys. Rev. A 73, 054302 (2006)


## Dirac Wave-Packet

## One-dimensional Dirac Hamiltonian:

$$
H_{\text {Dirac }}=\sigma_{z} p+\sigma_{x} m=\left(\begin{array}{cc}
p & m \\
m & -p
\end{array}\right), \quad \operatorname{eig}\left(H_{\text {Dirac }}\right)= \pm \sqrt{p^{2}+m^{2}}
$$

Dirac wave-packets

$$
\mathrm{a}=\text { localization parameter }
$$

$$
\begin{aligned}
\Psi_{ \pm}(x, t) & =N \int_{-\infty}^{+\infty} e^{-(a \pm i t) \sqrt{p^{2}+m^{2}}} e^{i p x}\left(I \pm \frac{H_{\text {Dirac }}}{\sqrt{p^{2}+m^{2}}}\right)\binom{1}{1} \\
& =N^{\prime}\binom{s_{ \pm}^{-1}(a \pm i(t+x)) K_{1}\left(m s_{ \pm}\right) \pm K_{0}\left(m s_{ \pm}\right)}{{s_{ \pm}^{-1}}^{-1}(a \pm i(t-x)) K_{1}\left(m s_{ \pm}\right) \pm K_{0}\left(m s_{ \pm}\right)}
\end{aligned}
$$

$$
s_{ \pm}=\sqrt{x^{2}+(a \pm i t)^{2}}
$$

Modified Bessel fxn

## DTQW Wave-Packet

## DTQW unitary mapping

$$
\begin{array}{ll}
U=e^{-i k \sigma_{z}} e^{-i \theta \sigma_{x}}=\left(\begin{array}{cc}
e^{-i k} \cos \theta & -i e^{-i k} \sin \theta \\
-i e^{i k} \sin \theta & e^{i k} \cos \theta
\end{array}\right), & e i g(U)=e^{ \pm i \omega(k)} \\
& \cos \omega(k)=\cos \theta \cos k
\end{array}
$$

DTQW "wave-packets" $\alpha$ = localization parameter

$$
\begin{aligned}
& \psi_{ \pm}(n, t)=N \int_{-\pi}^{+\pi} e^{-(\alpha \pm i t) \omega(k)} e^{i k n}\left(U-e^{ \pm i \omega(k)}\right)\binom{1}{1} d k \\
&=N^{\prime}\binom{I_{n}( \pm(t-1)-i \alpha)-e^{-i \theta} I_{n-1}( \pm t-i \alpha)}{I_{n}( \pm(t-1)-i \alpha)-e^{-i \theta} I_{n+1}( \pm t-i \alpha)} \\
& I_{n}(z)=\frac{1}{2 \pi} \int_{-\pi}^{+\pi} e^{-i z \omega(k)} e^{i k n} d k \approx i^{n} e^{-i z \pi / 2} J_{n}(z \cos \theta), \cos \theta \ll 1
\end{aligned}
$$

## Dirac

Relativistic Wave-Packet Spreading




Non-Relativistic Wave-Packet Spreading




Space

## Slices



## Dynamics governed by dispersion relations

$$
\begin{gathered}
\omega_{\text {DTQW }}=\cos ^{-1}(\cos \theta \cos p) \\
\omega_{\text {Dirac }}=\sqrt{p^{2}+m^{2}} \quad \omega_{\text {CTQW }}=2 \gamma(1-\cos p)
\end{gathered}
$$

Small $\theta$

```
                                    0=\pi/10
```




Each has $d \omega / d p<c=$ maximum speed

## Connecting the Quantum Walks

Introduce arbitrary coin parameter $\theta$, take various continuum limits!

$$
\begin{array}{|l}
\psi_{R}(n, \tau+1)=\cos \theta \psi_{R}(n-1, \tau)-i \sin \theta \psi_{L}(n-1, \tau) \\
\psi_{L}(n, \tau+1)=\cos \theta \psi_{L}(n+1, \tau)-i \sin \theta \psi_{R}(n+1, \tau)
\end{array}
$$



FWS, Phys. Rev. A 74, 030301 (R) (2006)
FWS, J. Math. Phys., 48, 082102 (2007)
Classical Analogy: Simple Diffusion vs. Telegrapher's Equation

## Connection for general graphs

- Childs (2008):
-Uses Szedegy's QW
-Hamiltonian simulation method

-QW "walks" on Hilbert space of dimension |G| ${ }^{2}$ (each state ~ edge of graph).
- D'Allesandro (2008):
-Controls QW by varying coin-flips.
-Lie theoretical proof of universal control
-QW on graphs of size |G|, coin flips of length log |G|


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## Algorithms

## - NAND Tree:



Properties of graph revealed by scattering Farhi, Goldstone, \& Guttman (2008)

- Graph Isomorphism

Properties of graph revealed by multiple walkers (Wang + others)


- Universality

ALL Q Computations can be written as a CTQW Childs (2008)

## Encodings and Implementations

- Quantum Computer
-Graph vertices encoded in quantum bits
-Walks encoded by unitary logic gates
-Most efficient encoding has $\log _{2}|\mathrm{G}|$ qubits $=$ binary encoding
- Physical Network
-Graph vertices encoded in degrees of freedom (atom sites, electron sites, etc.)
-Walks encoded by natural Hamiltonian
-Most natural encoding has |G| qubits = unary encoding


## QW Implementations

(not an exhaustive list)

| Proposal | System | Walk | Encoding |
| :---: | :---: | :---: | :---: |
| W. Dür et al. (2002) | Optical Lattices | ID DTQW | Unary |
| B. C. Sanders et al. (2002) | Cavity-QED | ID DTQW | Unary/Binary |
| A. Romito et al. (2005) | Superconductors | ID CTQW | Unary |
| C.A. Ryan et al. (2005) | NMR exp | 2D DTQW | Binary |
| J. M.Taylor (2007) | Quantum Dots | NAND CTQW | Unary |
| F. w. Strauch (2008) | Superconductors | Hypercube | Unary |
| Karski et al. (2009) | Optical Lattices | ID DTQW | Unary |
| Zähringer et al. (2010) | Ion Traps | ID DTQW | Unary |
| Broome et al. (2010) | Photons | ID DTSQW | Unary |

Most proposals encode the walk using a unary representation.
Does this really demonstrate quantum computational speedup?
DTQW = Discrete-time quantum walk CTQW = Continuous-time quantum walk

## Nature's Implementation? Photosynthesis

FMO Complex


LH2 Complex


Unary QW with disorder \& decoherence
M. Mohseni, P. Rebentrost, S. Lloyd, and A.Aspuru-Guzik,
J. Chem. Phys. I 29, I74I06 (2008)

## Decoherence in the Hypercube Quantum Walk

Two "natural" models of decoherence for the hypercube:

- Subspace Model
- Each node is represented by a state of several qubits.
- Binary encoding with d qubits
- Dephasing occurs between subspaces of the hypercube.
- Useful model for quantum computer implementation of quantum walks.
G.Alagic and A. Russell,

Phys. Rev.A 72, 062304 (2005)

- Vertex Model
- Each node is represented by the excited state of one qubit.

Unary encoding with $\mathbf{2 d}^{\mathbf{d}}$ qubits

- Dephasing occurs between vertices of the hypercube.
- Useful model for quantum state transfer in qubit networks.


## Decoherence in the Hypercube.

## Quantum Walk

## Lindblad Equation

$$
\frac{d \rho}{d t}=-i[H, \rho]-\sum_{j} \lambda_{j}\left(L_{j} \rho L_{j}^{+}-\frac{1}{2} L_{j}^{+} L_{j} \rho-\frac{1}{2} \rho L_{j}^{+} L_{j}\right)
$$

Dephasing leads to decay of off-diagonal elements of $\rho$

## Subspace Model ( $x=x_{1} x_{2} x_{3} x_{1} x_{5 . .}$ )

$$
\frac{d \rho_{x, y}}{d t}=-i \sum_{z}\left(H_{x z} \rho_{z, y}-\rho_{x, z} H_{z y}\right)-\lambda \sum_{i}\left(1-\delta_{x, y,}\right) \rho_{x, y}
$$

Dephasing between different subspaces $\left(x_{j} \neq y_{i}\right)$

## Vertex Model

$$
\frac{d \rho_{x, y}}{d t}=-i \sum_{z}\left(H_{x z} \rho_{z, y}-\rho_{x, z} H_{z y}\right)-\lambda\left(1-\delta_{x y}\right) \rho_{x, y}
$$

Dephasing between different vertices ( $x \neq y$ )

## Comparison of Models



See F. W. Strauch, Phys. Rev.A 79, 0323 I9 (2009) for details


Decoherence in the vertex model has a hitting probability with a lower bound independent of the hypercube dimension $\boldsymbol{d}$ ! Quantum speedup robust for unary encoding.

## Quantum State Transfer

One can transfer the state of a single qubit from site $A$ to site $B$ using a set of permanently coupled qubits with Hamiltonian:

$$
H=-\frac{1}{2} \sum_{j} \hbar \omega_{j} \sigma_{j}^{z}+\sum_{j k} \hbar \Omega_{j k}\left(\sigma_{j}^{+} \sigma_{k}^{-}+\sigma_{k}^{+} \sigma_{j}^{-}\right)
$$

Dynamics of a single excitation (with $\omega=0$ ) maps onto the continuous-time quantum walk with $H=\hbar \Omega$, where the coupling matrix $\Omega$ is proportional to the adjacency matrix of the coupling graph. Certain choices of couplings such as the hypercube lead to perfect state transfer:


## Hypercube State Transfer

- Each vertex represents a qubit. Quantum states travel along all paths simultaneously in superposition with full constructive interference, yielding perfect state transfer.



## Superconducting Circuits

Artificial Atoms $\rightarrow$ Artificial Molecules $\rightarrow$ Artificial Solids


Qubits + Resonators


Phase Qubit


UC Santa Barbara


## Perfect State Transfer with Phase

 Qubits- Programmable: Any two nodes can communicate by programming the qubit frequencies in the network.
- Parallel: Multiple quantum states can be transferred at the same time.
- Efficient: Transfer time is independent of the distance between nodes!
- High Fidelity: $\mathrm{F}>90 \%$ possible using existing technology, modest dimensions! Requires Study of Disorder and Decoherencen FWS and C.J.Williams, Phys. Rev. B 78, 0945 I 6 (2008)

C. Chudzicki and FWS, Phys. Rev. Lett. I05, 26050I (20I0)


## Conclusions

- Quantum Walks are an exciting testbed for quantum information processing
- Connections have emerged between and but each has their own advantages (CTQW for studying physical networks, DTQW for computer algorithms)
- Search algorithms well developed---new algorithms on the way?
- Implementations require study of encoding, decoherence, and disorder. Understanding robustness may be the key.

